

High-Fidelity Specification of Real-World Devices

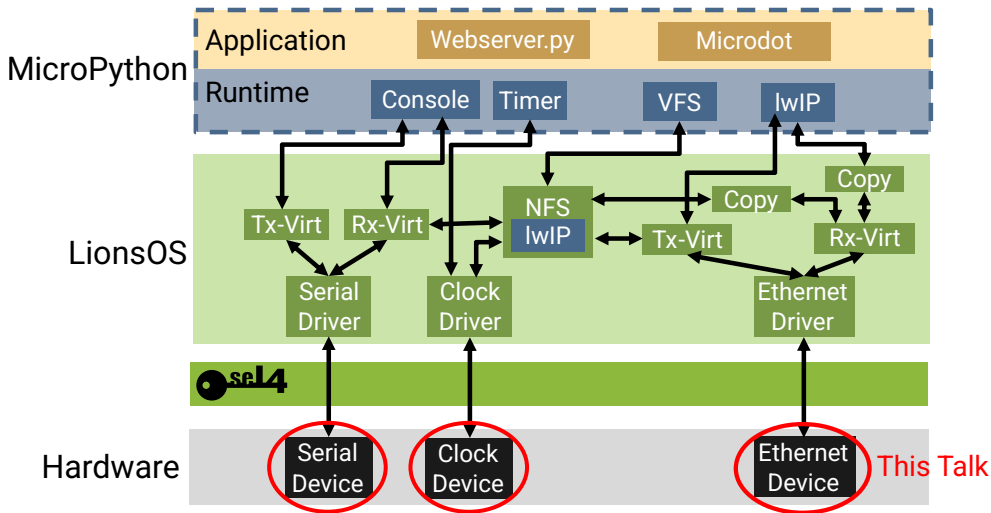
Liam Murphy¹ Albert Rizaldi² Lesley Rossouw¹
Chen George³ James Treloar¹ Hammond Pearce¹
Miki Tanaka¹ Gernot Heiser¹

¹UNSW Sydney ²PlanV GmbH ³University of Wisconsin - Madison

PLOS '25

Background

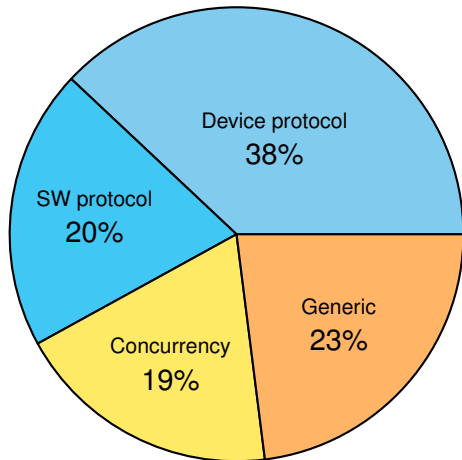
Who Are We?



Why Drivers?

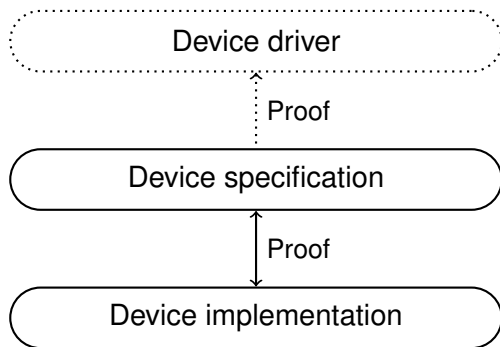


- Device driver bugs are a major source of OS vulnerabilities
 - Cause of majority of Linux CVEs from 2018-2022 [Pohjola et al. 2023]
- Dominant cause of Linux driver bugs is device-protocol violations [Ryzhyk et al. 2009]
 - Without knowledge of the device protocol, formal verification cannot prevent these



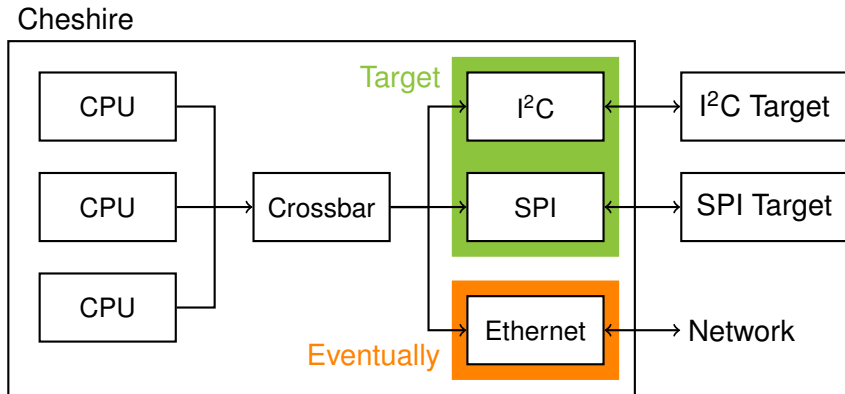
Causes of Linux driver bugs (2002–2008)

- Prevent device-protocol bugs by:
 - Formally specifying device interfaces
 - Verifying these specifications against the device implementation
- Future work: driver verification



- Target hardware
- Verification approach
- Specification format
- Application
- Status

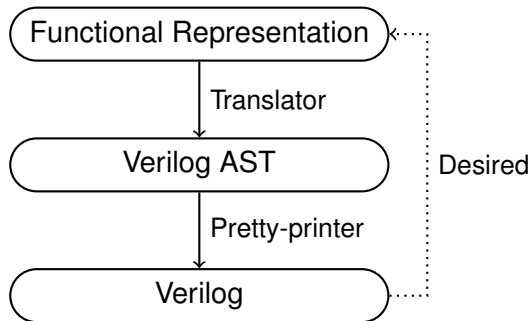
Target Hardware



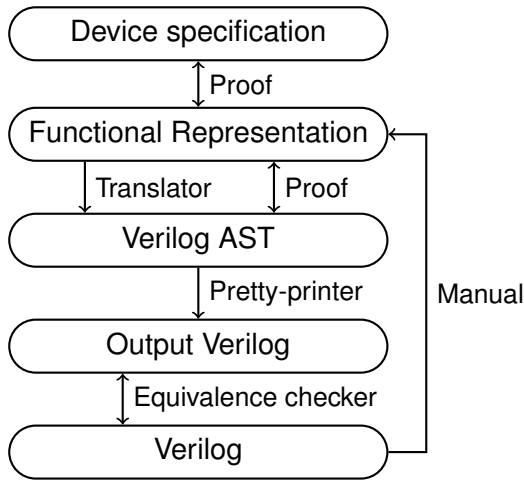
	Verilog	Driver	HOL spec
I²C	5993	713	1414
SPI	4609	864	1235
Ethernet	3987	641	N/A

Verification Approach

- HOL4 theorem prover
- HOL4 Verilog formalisation [Löow and Myreen 2019]:
 - Semantics for Verilog AST
 - Semantics for functional representation
 - Proof-producing translator from functional to AST
 - Pretty-printer for deep embedding
- Problem: need opposite direction

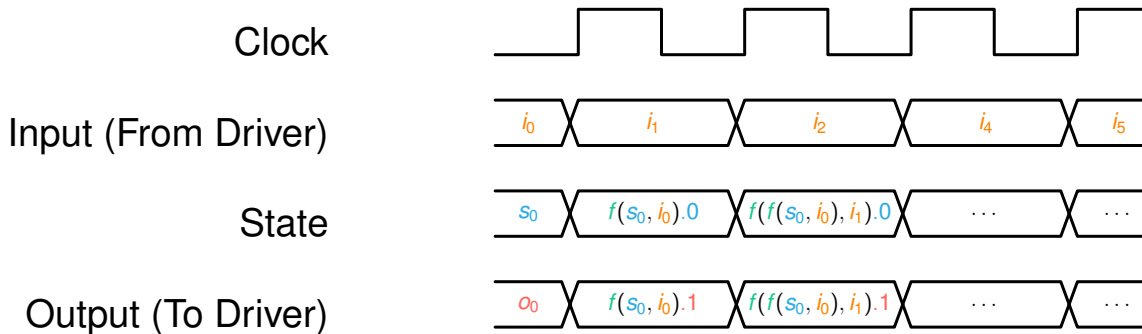


- Workaround:
 - Manually translate Verilog to functional representation
 - Translate and export into Verilog
 - Use equivalence-checker (eqy) to prove equivalence to original Verilog
- Finally, prove that functional representation refines specification



Specification Format

Specification Format

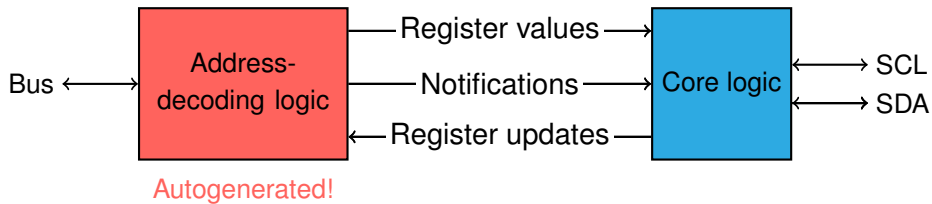


$$f: s \rightarrow i \rightarrow s \times o$$

- Model can be non-deterministic
 - Inputs not controlled by driver
 - Underspecification
- f only returns one possible state
- Determined by `fnums` - infinite stream of numbers
- To obtain set of all possible states, try all `fnums`

Application to Cheshire

Structure of Cheshire Peripherals



- Format of Cheshire peripheral specifications:

```
cheshire_run tick read write:
```

```
state -> req option list ->
```

```
ffi_outcome + state # word32 option list
```

- Input instantiated with optional MMIO request
- Output instantiated with response to read requests
- Executes multiple cycles

```
cheshire_run tick read write:  
  state -> req option list ->  
  ffi_outcome + state # word32 option list
```

- tick represents core logic
- read, write and cheshire_run represent address-decoding logic

Status




	I ² C		SPI	
	Address-decoding	Core	Address-decoding	Core
Specification	✓	✓	✓	✓
Functional Representation	✓	✓	✓	WIP
Equivalence-checking	✓	WIP	✓	WIP
Correctness proof	✓	Almost!	✓	WIP

Questions?



Note

UNSW is recruiting OS faculty members - talk to Gernot if you're interested

-  Pohjola, Johannes Åman et al. (Oct. 2023). “Pancake: Verified Systems Programming Made Sweeter”. In: *Workshop on Programming Languages and Operating Systems (PLOS)*. Koblenz, DE.
-  Ryzhyk, Leonid et al. (Apr. 2009). “Dingo: Taming Device Drivers”. In: *EuroSys Conference*. Nuremberg, DE, pp. 275–288.
-  Löow, Andreas and Magnus O. Myreen (2019). “A proof-producing translator for Verilog development in HOL”. In: *Proceedings of the International Workshop on Formal Methods in Software Engineering (FormaliSE@ICSE)*, pp. 99–108.